# Many-Core Virtual Machines Decoupling Abstract From Concrete Concurrency





raceMonkey Bytecode

AST

Bvtecode





stack

stack

0 switch, threaded

1threaded,

variable 1 or 2

variable >= 1

71

234

## concurrency support is limited

- VM support is minimal
  - only one specific concurrency model is supported
  - only few ILs provide notion of concurrency
  - no comprehensive abstraction

<b></b>	NET 2	inferno erma	ERLANG	molart
Model	Threads/Locks	CSP	Actors	Data-flow
IL Support	Marginal	High-level	High-level	Marginal
StdLib	Low/high-level	High-level	High-level	High-level



Stefan Marr, Michael Haupt, and Theo D'Hondt Intermediate Language Design of High-level Language Virtual Machines: owards Comprehensive Concurrency Support

# a VM has to:

## decouple abstract concurrency models

- abstract concurrency models are defined by languages or libraries
- used by application developers



• wide range of models supported by VM is necessary • implementing unsupported models on top is hard restrictions hinder efficient implementation • support at VM-level allows reuse and optimization

and concrete concurrency models

concrete concurrency models are provided by the underlying system



#### Single-Core

- preemptive OS thread
- instruction-level
- parallelism
- VM challenges
- deep cache hierarchies • cache-consciousness
- required



Main M	Memory

#### Multi-Core

- uniform memory acces
- native support for thread-level parallelism
- and cache coherency locality and cache hierar-
- chy must be considered • avoid cache thrashing





### Many-Core

- non-uniform memory access architectures
- can have explicit coreto-core communication
- very diverse designs • with/out cache coher.
- explicit inter-core com.





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# what to include in ILs?

there are various ways to express concurrency and solutions are domain-specific



# how to combine the various models?



LOAD SPAWN? WAIT? LOCK? COMMIT? POP PUSH









Is the notion of locality inevitable for a VM?

Non-shared Memory Language natural fit for NUMA



Partitioned Global Address Space locality explicit in shared-memory Non-Uniform **Memory Access** 



d in

protecting and isolating shared state



⊒. avoiding mutable shared state perhaps allowing immutable shared state

Can a VM support this better by some notion of Encapsulation?



http://soft.vub.ac.be/~smarr/research/



